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01 Background

# Background



At the Seventh International World Wide Web conference(WWW98), Sergey Brin and Larry Page's paper "**The PageRank citation ranking: Bringing order to the Web**" made small ripples in the information science community that quickly turned into waves.



#### april 14 to 18, 1998

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reto ambühler (www.ethz.ch/r\_amb/)





## Great Success 🔶 Hyperlink Structure 🖕 PageRank



# Markoy Chain

## Markov Chain



#### **Markov Chain:**





**Irreducible**(不可约): any state can reached from any state.

Absorbing states(吸收态): the probability of leaving this state is zero



## Markov Chain:

Limit Theorem(极限定理):

A homogeneous(齐次), irreducible(不可约), aperiodic(非周期) and positive recurrent(正常返) Markoc Chain has:

a limiting distribution:  $\lim_{n \to \infty} p_{ij}^{(n)} = \pi_j, i, j \in E$ 

also a stationary distribution:  $\pi P = \pi$   $\pi_j \ge 0, \sum_{j \in E} \pi_j = 1$ 



### Markov Chain:

#### Limit Theorem(极限定理):

$$P^{(n)} = (p_{ij}^{(n)}) = \begin{bmatrix} p_{11}^{(n)} & p_{12}^{(n)} & \cdots & p_{1j}^{(n)} & \cdots \\ p_{21}^{(n)} & p_{22}^{(n)} & \cdots & p_{2j}^{(n)} & \cdots \\ p_{31}^{(n)} & p_{32}^{(n)} & \cdots & p_{3j}^{(n)} & \cdots \\ \vdots & \vdots & & \vdots \end{bmatrix} = \begin{bmatrix} \pi_1 & \pi_2 & \cdots & \pi_j & \cdots \\ \pi_1 & \pi_2 & \cdots & \pi_j & \cdots \\ \pi_1 & \pi_2 & \cdots & \pi_j & \cdots \\ \vdots & \vdots & & \vdots \end{bmatrix} = \begin{bmatrix} \Pi \\ \Pi \\ \Pi \\ \vdots \end{bmatrix}, n \to \infty$$

# Markov Chain



#### **Markov Chain(Aperiodic)**:

**Period** of a state:

 $k = g.c.d\{n > 0: \Pr(X_n = i | X_0 = i) > 0\}$ 

divisor : the greatest common

**Aperiodic(**非周期): *k* = 1



#### **Markov Chain(Positive Recurrent)**:

The first return probability(首返概率):

$$f_{ii}^{(n)} = P\{X_n = i, X_k \neq i, 1 \le k < n | X_0 = i\}$$

Recurrent (常返): 
$$f_{ii} = \sum_{n=1}^{\infty} f_{ii}^{(n)} = 1$$

Positive Recurrent (正常返): 
$$\mu_i = \sum_{n=1}^{\infty} n f_{ii}^{(n)} < +\infty$$

 $\infty$ 



#### **Markov Chain**:

#### Positive Recurrent (正常返):



**State 1:**  $f_{11}^{(1)} = 1, f_{11}^{(n)} = 0 (n > 1) \Rightarrow f_{11} = 1, \mu_1 = 1$  *positive recurrent !* 

State 2: 
$$f_{22}^{(1)} = 0, f_{22}^{(2)} = pq, f_{22}^{(n)} = 0 (n > 2) \Rightarrow f_{22} = \sum_{n=1}^{\infty} n f_{22}^{(n)} = pq < 1$$
  
*transient* !

# 03 The Basic PageRank Model

# The Basic PageRank Model

## Page Links:



## If a page

- is linked by many pages;
- which link to it is authoritative

#### Then

it will gain high PageRank!







#### Markov matrix/stochastic matrix:

- Every element is Nonnegative
- Row/Column sums are +1

#### **Properties:**

Spectral radius(the supremum among the absolute values of spectrums) is 1;





If every element is positive, **Markov matrix** is irreducible (strongly connected):

# Every irreducible markov matrix has a stationary vector:

 $\pi$ : does not change under application of the transition matrix:

$$\pi^T P = \pi^T$$

It's independent of the initial state

PageRank vector is  $\pi$ , so this is just an eigenvector problem!



The independence of the initial state:

$$u_{0} = (0.02, 0.98)$$

$$P = \begin{bmatrix} 0.2 & 0.05 \\ 0.8 & 0.95 \end{bmatrix}$$

$$Px = \lambda x$$

$$u_{0} = \begin{bmatrix} 0.02 \\ 0.8 \end{bmatrix} = \begin{bmatrix} 0.2 \\ 1 \end{bmatrix} = 0.75 \begin{bmatrix} -1 \\ 1 \end{bmatrix}$$

$$u_{0} = \begin{bmatrix} 0.02 \\ 0.98 \end{bmatrix} = \begin{bmatrix} 0.2 \\ 0.8 \end{bmatrix} + 0.18 \begin{bmatrix} -1 \\ 1 \end{bmatrix}$$

$$u_{k} = 1 \begin{bmatrix} 0.2 \\ 0.8 \end{bmatrix} + 0.75^{k} * 0.18 \begin{bmatrix} -1 \\ 1 \end{bmatrix} \bigstar$$

$$u_{0} = 1 \begin{bmatrix} 0.2 \\ 0.8 \end{bmatrix} + 0.75 * 0.18 \begin{bmatrix} -1 \\ 1 \end{bmatrix}$$



Practical web graphs are not necessarily strongly connected.

• For Dangling Nodes:

$$\bar{P} = P + av^T$$

a:  $a_i = 1$  if row *i* of **P** corresponds to a dangling node, and 0, otherwise.  $v^T$ : a general probability vector.

• For Irreducible:

$$\overline{\bar{P}} = \alpha \overline{P} + (1 - \alpha) e v^T$$

 $e = [1, 1, 1, \dots, 1]^T$ 



$$P = \begin{pmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix} \Rightarrow \bar{P} = \begin{pmatrix} 0 & 1/2 & 1/2 & 0 & 0 & 0 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ 1/3 & 1/3 & 0 & 0 & 1/3 & 0 \\ 0 & 0 & 0 & 0 & 1/2 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1/2 & 0 & 1/2 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{pmatrix}$$
$$\bar{P} = \begin{pmatrix} 1/60 & 7/15 & 7/15 & 1/60 & 1/60 & 1/60 \\ 1/6 & 1/6 & 1/6 & 1/6 & 1/6 & 1/6 \\ 19/60 & 19/60 & 1/60 & 1/60 & 19/60 & 1/60 \\ 1/60 & 1/60 & 1/60 & 1/60 & 7/15 & 7/15 \\ 1/60 & 1/60 & 1/60 & 1/60 & 7/15 & 1/60 & 7/15 \\ 1/60 & 1/60 & 1/60 & 11/12 & 1/60 & 1/60 \end{pmatrix} \alpha = 0.9$$
$$v = \begin{bmatrix} \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6} \end{bmatrix}^{T}$$



$$\bar{\bar{P}} = \alpha \bar{P} + (1 - \alpha) e v^T$$

Another Interpretation for  $\overline{\overline{P}}$  is:

Even though the customer always browse webpages by hyperlinks, but he can also use URL to "teleport" to a new page.



# 04 The Power Method



The Power Method:

$$\mathbf{x}^{(k)T} = \mathbf{x}^{(k-1)T} \mathbf{\bar{\bar{P}}} = \alpha \mathbf{x}^{(k-1)T} \mathbf{\bar{P}} + (1-\alpha) \mathbf{x}^{(k-1)T} \mathbf{ev}^T$$
$$= \alpha \mathbf{x}^{(k-1)T} \mathbf{\bar{P}} + (1-\alpha) \mathbf{v}^T$$
$$= \alpha \mathbf{x}^{(k-1)T} \mathbf{P} + (\alpha \mathbf{x}^{(k-1)T} \mathbf{a} + (1-\alpha)) \mathbf{v}^T,$$

 $\mathbf{x}^{(k-1)T}e = \mathbf{1}.$ 



Four advantages:

- $\overline{\overline{P}}$  and  $\overline{P}$  are never formed or stored.
- Since P is sparse, each vector-matrix multiplication can be computed in nnz(P) flops.
   nnz(P): the number of nonzeros in P.
- At each iteration, the power method only requires the storage of one vector, the current iterate.
- Converges quickly.

# 05 Discussion about the Model



 $\alpha$ (damping factor):

Google always using = 0.85, so why this choice for ?

1. a trade-off:

- The larger is, the more the true hyperlink structure of the web is used to determine webpage importance.
- The smaller is, the faster the convergence for power method.



α:

Rate of convergence(r) is determined by the subdominant eigenvalue( $\lambda_2$ ) of the transition matrix  $\overline{\overline{p}}$ , and this eigenvalue is strictly less than  $\alpha$  for an irreducible hyperlink matrix.

 $r \propto (\lambda_1 - \lambda_2) = (1 - \lambda_2) \propto (1 - \alpha)$ 

for  $\tau = x^{(k+1)T} - x^{(k)T} = 10^{-8}$ , only need 142 iterations.



α:

- 2. intuitive reality:
- $\alpha$  = 0.85 implies that roughly:
- 5/6 of the time a Web surfer randomly clicks on hyperlinks.
- 1/6 of the time this Web surfer will go to the URL line and type the address of a new page.



 $\mathbf{v}^{\mathbf{T}}$ :

Instead of traditional  $v^T = \frac{1}{n}e^T$ ,  $v^T$  can be more personalized.

- different classes of surfers have different  $v^T$ .
- control spamming done by the so-called link farms.



## **Discussion about the Model**



	Т		
V	•	•	
v		•	

	$\begin{array}{c} \text{Damping} \\ \text{Factor} \\ (\alpha) \end{array}$	Personalization Vector (v)	$\begin{array}{c} \text{Google} \\ \text{Matrix} \\ (G) \end{array}$	PageRank Vector $(\approx \pi)$	Ordering of Nodes (1 = Highest)
Model 1	0.85	$\begin{pmatrix} \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} \end{pmatrix}$	$\begin{pmatrix} \frac{3}{80} & \frac{71}{80} & \frac{3}{80} & \frac{3}{80} \\ \frac{3}{80} & \frac{3}{80} & \frac{71}{80} & \frac{3}{80} \\ \frac{37}{80} & \frac{3}{80} & \frac{3}{80} & \frac{37}{80} \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} \end{pmatrix}$	(0.21 0.26 0.31 0.21)	(3 2 1 3)
Model 2	0.85	(1 0 0 0)	$\begin{pmatrix} \frac{3}{20} & \frac{17}{20} & 0 & 0\\ \frac{3}{20} & 0 & \frac{17}{20} & 0\\ \frac{23}{40} & 0 & 0 & \frac{17}{40}\\ \frac{29}{80} & \frac{17}{80} & \frac{17}{80} & \frac{17}{80} \end{pmatrix}$	(0.30 0.28 0.27 0.15)	(1 2 3 4)
Model 3	0.95	$\begin{pmatrix} \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} \end{pmatrix}$	$\begin{pmatrix} \frac{1}{80} & \frac{77}{80} & \frac{1}{80} & \frac{1}{80} \\ \frac{1}{80} & \frac{1}{80} & \frac{77}{80} & \frac{1}{80} \\ \frac{39}{80} & \frac{1}{80} & \frac{1}{80} & \frac{39}{80} \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} \end{pmatrix}$	(0.21 0.26 0.31 0.21)	(3 2 1 3)
Model 4	0.95	(1 0 0 0)	$\begin{pmatrix} \frac{1}{20} & \frac{19}{20} & 0 & 0\\ \frac{1}{20} & 0 & \frac{19}{20} & 0\\ \frac{21}{40} & 0 & 0 & \frac{19}{40}\\ \frac{23}{80} & \frac{19}{80} & \frac{19}{80} & \frac{19}{80} \end{pmatrix}$	(0.24 0.27 0.30 0.19)	(3 2 1 4)



# Forcing Irreducibility:

enforce every node is directly connected to every other node. (alter the true nature of the Web)

add a dummy node to the Web which connects to every other node and to which every other node is connected.

$$\hat{\mathbf{P}} = \begin{pmatrix} \alpha \bar{\mathbf{P}} & (1-\alpha)\mathbf{e} \\ \frac{\mathbf{v}^T}{\mathbf{v}^T} & 0 \end{pmatrix}$$

# Forcing Irreducibility:

## LeaderRank:

$$s_i(t+1) = \sum_{j=1}^{N+1} \frac{a_{ji}}{k_j^{out}} s_j(t)$$
$$S_i = s_i(t_c) + \frac{s_g(t_c)}{N}$$

- self-adaptive
- parameter-free





06 Other Topics





Storage and Speed

"The World's Largest Matrix Computation"

- Spam
- The evolution and dynamics of the Web
- Web's structure:

how to use the *scale-free* structure to improve PageRank computations

#### • Community:

how do changes within the community affect the PageRank of community pages?



